

## A NEW APPROACH TO SOLVE MINIMUM SPANNING TREE PROBLEM: MAXIMUM COST PRUNING METHOD

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**Abstract:** The Minimum Spanning Tree for a given graph is the Spanning Tree of minimum cost for that graph. There are various algorithms to solve Minimum Spanning Tree problem. Existing two algorithms are: Prim's algorithm and Kruskal's algorithm. Time complexity of Prim's and Kruskal's algorithm is  $O(E \lg V)$ . Most of algorithms for Minimum Spanning Tree work only for undirected graph. Our ongoing research will focus on Minimum Spanning Tree problem for both directed and undirected graph.  
**Index Terms:** Spanning tree, Minimum Spanning Tree.

### I. INTRODUCTION

A graph G consists of V and E. V is set vertex and E is edge that connecting two vertex. There are basic two types of graph: **Undirected graph and Directed graph**. Graph can be represented in two ways: Adjacency Matrix and Adjacency List.

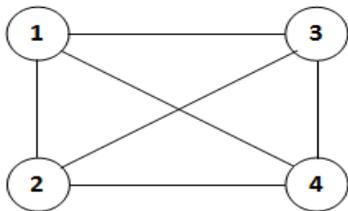
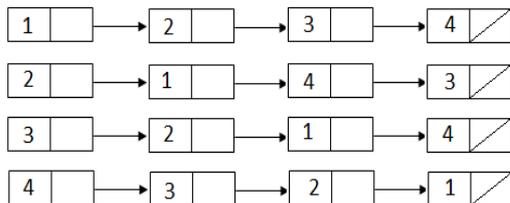


Fig.1.1 [Graph]

This graph can be represented using Adjacency List as follows:



This graph can be represented using Adjacency Matrix as follows:

$$\begin{matrix}
 & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\
 \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix}
 0 & 1 & 1 & 1 \\
 1 & 0 & 1 & 1 \\
 1 & 1 & 0 & 1 \\
 1 & 1 & 1 & 0
 \end{bmatrix}
 \end{matrix}$$

Spanning tree of the graph is a tree that contains all the node and doesn't contain cycle.

A **Minimum Spanning Tree (MST)** of a weighted graph is a Spanning Tree in which the sum of the weights of all its edges is minimum of all such possible spanning trees of the graph. The main aim of the Minimum Spanning Tree algorithm is to find the shortest path for given graph in which all nodes will be visited exactly once.

### II. MINIMUM SPANNING TREE ALGORITHM

#### A. Prim's algorithm

In prim's algorithm starting node is selected randomly from given graph. A set which contain all the edges in the graph is created. The edge, which is adjacent to the randomly selected node and with minimum weight among all adjacent edge, is selected from the set and add to new graph if it connects a vertex in the tree with a vertex not in the tree. This process is repeated until every edge in the set connects two vertices in the tree. Time complexity of this algorithm is  $O(E \lg V)$

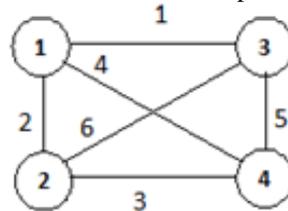


Fig.2.1 [Initial Graph]



Fig.2.2 [Step-1]

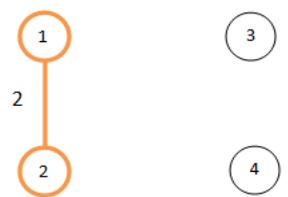


Fig.2.3 [Step-2]

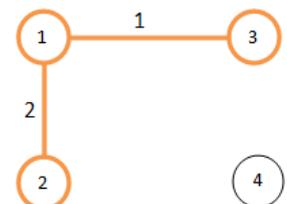


Fig.2.4 [Step-3]

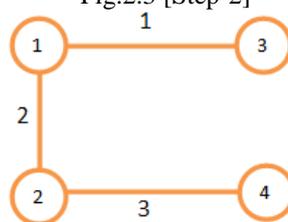


Fig.2.5 [Step-4]

Total Cost of the MST:  $2+1+3 = 6$



```
nirav@nirav: ~/Desktop
deleted edge is from 8 to 2 with weight:20
updated matrix is:
0 10 0 0 20 20 5 15 0 0 0
0 0 5 10 0 0 0 0 0 0 0
0 15 0 5 0 0 0 0 0 0 0
0 0 0 0 10 0 0 0 0 0 0
0 0 0 0 0 5 0 0 0 0 0
0 0 0 0 0 0 10 0 0 0 0
5 0 0 0 0 0 5 0 0 0 0
0 15 0 0 0 0 0 0 20 0 10
0 5 15 0 0 0 0 0 0 0 0
indegree for node 1 is:1
indegree for node 2 is:4
indegree for node 3 is:2
indegree for node 4 is:3
indegree for node 5 is:2
indegree for node 6 is:3
indegree for node 7 is:2
indegree for node 8 is:1
indegree for node 9 is:0
indegree for node 10 is:1
no.of edges=19
```

```
nirav@nirav: ~/Desktop
Minimum Spanning Tree is:
0 9 0 0 0 0 0 0 0 0 0
9 0 0 0 0 3 0 0 0 0 0
0 0 0 0 7 0 0 0 0 0 0
0 0 0 0 5 0 0 0 0 0 0
0 0 7 5 0 9 0 0 0 0 0
0 3 0 0 9 0 11 4 0 0 0
0 0 0 0 0 11 0 0 0 0 0
0 0 0 0 0 4 0 0 0 6 0
0 0 0 0 0 0 0 6 0 2
0 0 0 0 0 0 0 0 2 0
Total time = 0.003284 seconds
nirav@nirav:~/Desktop$
```

```
nirav@nirav: ~/Desktop
minimum spanning tree is:
0 0 0 0 0 5 0 0 0 0 0
0 0 5 0 0 0 0 0 0 0 0
0 0 0 5 0 0 0 0 0 0 0
0 0 0 0 10 0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0
5 0 0 0 0 0 5 0 0 0 0
0 0 0 0 0 0 0 20 0 10
0 5 0 0 0 0 0 0 0 0 0
Total time = 0.008343 seconds
nirav@nirav:~/Desktop$
```

V. CONCLUSION

Here we can conclude that existing algorithm (Prim's and Kruskal's algorithm) cannot give proper solution for Directed Graph. Proposed algorithm (Maximum Cost Pruning method) gives proper solution for both directed and undirected graph.

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- Existing Prim's and Kruskal's algorithm can't give solution for directed graph

B. Undirected Graph

- Weight matrix for undirected graph

```
nirav@nirav: ~/Desktop
source node destination node weight
1 2 9
1 6 12
1 7 20
2 3 10
2 4 3
3 4 8
3 5 7
4 5 5
4 10 21
5 6 9
5 9 14
5 10 18
6 8 4
7 8 11
7 13 13
8 9 6
9 10 2
weight matrix:
0 9 0 0 0 12 20 0 0 0 0
9 0 10 0 0 7 0 0 0 0 0
0 10 0 8 0 7 0 0 0 0 0
0 0 8 0 5 0 0 0 0 0 21
0 0 0 7 5 0 9 0 0 14 18
12 3 0 0 0 9 0 11 4 0 0
20 0 0 0 0 0 4 0 13 0 6
0 0 0 0 0 0 4 0 13 0 6
0 0 0 0 0 14 0 0 6 0 2
0 0 0 0 21 18 0 0 0 2 0
no.of edges=17
```

- Solution using Maximum Cost Pruning method

```
nirav@nirav: ~/Desktop
edge from 7 to 6 with weight 11 can't deleted
no.of edges=11
edge from 3 to 2 with weight 10 is deleted
updated matrix is:
0 9 0 0 0 0 0 0 0 0 0
9 0 0 0 0 3 0 0 0 0 0
0 0 0 8 7 0 0 0 0 0 0
0 0 8 0 5 0 0 0 0 0 0
0 0 7 5 0 9 0 0 0 0 0
0 3 0 0 9 0 11 4 0 0 0
0 0 0 0 0 11 0 0 0 0 0
0 0 0 0 0 4 0 0 6 0 0
0 0 0 0 0 0 0 6 0 2
0 0 0 0 0 0 0 0 2 0
no.of edges=10
edge from 6 to 5 with weight 9 can't deleted
no.of edges=10
edge from 2 to 1 with weight 9 can't deleted
no.of edges=10
edge from 4 to 3 with weight 8 is deleted
```